

JAWAHARLAL NEHRU TECHNOLOGICAL UNIVERSITY HYDERABAD

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Lecture-21 By Vivek Dubey

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COMPUTER SCIENCE AND ENGINEERING

CS403PC

COMPUTER SCIENCE AND ENGINEERING (ARTIFICIAL INTELLIGENCE AND MACHINE LEARNING)

CS403PC

BHARAT INSTITUTE OF ENGINEERING AND TECHNOLOGY

Scheme of Operating System

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S. No.	Course Code	Course Title	L	T	P	Credits
3	CS403PC	Operating Systems	3	0	0	3
6	CS406PC	Operating Systems Lab	0	0	3	1.5

Aim

Program to implement Deadlock Avoidance & Prevention by using Banker's Algorithm

Concept of Banker's Algorithm

- The banker's algorithm is a resource allocation and deadlock avoidance algorithm
- It tests for safety by simulating the allocation for predetermined maximum possible amounts of all resources,
- Also, it then makes an “s-state” check to test for possible activities, before deciding whether allocation should be allowed to continue.

Why Banker's algorithm is named so?



- Banker's algorithm is named so because it is used in banking system to check whether loan can be sanctioned to a person or not.
- Suppose there are n number of account holders in a bank and the total sum of their money is S .
- If a person applies for a loan then the bank first subtracts the loan amount from the total money that bank has and if the remaining amount is greater than S then only the loan is sanctioned.
- It is done because if all the account holders comes to withdraw their money then the bank can easily do it.
- In other words, the bank would never allocate its money in such a way that it can no longer satisfy the needs of all its customers. The bank would try to be in safe state always.

Data structures to implement the Banker's Algorithm

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- Let ' n ' be the number of processes in the system and ' m ' be the number of resources types.

Available :

- It is a 1-d array of size ' m ' indicating the number of available resources of each type.
- $\text{Available}[j] = k$ means there are ' k ' instances of resource type R_j
- **Max :**
 - It is a 2-d array of size ' $n*m$ ' that defines the maximum demand of each process in a system.
 - $\text{Max}[i, j] = k$ means process P_i may request at most ' k ' instances of resource type R_j .



- **Allocation :**

- It is a 2-d array of size ' $n*m$ ' that defines the number of resources of each type currently allocated to each process.
- $\text{Allocation}[i, j] = k$ means process P_i is currently allocated ' k ' instances of resource type R_j

- **Need :**

- It is a 2-d array of size ' $n*m$ ' that indicates the remaining resource need of each process.
- $\text{Need}[i, j] = k$ means process P_i currently need ' k ' instances of resource type R_j
- $\text{Need}[i, j] = \text{Max}[i, j] - \text{Allocation}[i, j]$



Example:

Considering a system with five processes P_0 through P_4 and three resources of type A, B, C.

Resource type A has 10 instances, B has 5 instances and type C has 7 instances.

Suppose at time t_0 following snapshot of the system has been taken.

Considering a system with five processes P_0 through P_4 and three resources of type A, B, C

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Resource type A has 10 instances, B has 5 instances and type C has 7 instances.



Process	Allocation	Max	Available
	A B C	A B C	A B C
P_0	0 1 0	7 5 3	3 3 2
P_1	2 0 0	3 2 2	
P_2	3 0 2	9 0 2	
P_3	2 1 1	2 2 2	
P_4	0 0 2	4 3 3	

Question1. What will be the content of the Need matrix?

$\text{Need}[i, j] = \text{Max}[i, j] - \text{Allocation}[i, j]$

So, the content of Need Matrix is:

Process	Need		
	A	B	C
P_0	7	4	3
P_1	1	2	2
P_2	6	0	0
P_3	0	1	1
P_4	4	3	1

Algorithm

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- Step 1: Start
- Step 2: Get the values of resources and processes.
- Step 3: Get the avail value.
- Step 4: After allocation find the need value.
- Step 5: Check whether its possible to allocate.
- Step 6: If it is possible then the system is in safe state.
- Step 7: Else system is not in safety state.
- Step 8: If the new request comes then check that the system is in safety or not if we allow request.
- Step 9: Stop

Run

INPUT

Enter number of processes	--	5		
Enter number of resources	--	3		
Enter details for P0				
Enter allocation	--	0	1	0
Enter Max	--	7	5	3
Enter details for P1				
Enter allocation	--	2	0	0
Enter Max	--	3	2	2
Enter details for P2				
Enter allocation	--	3	0	2
Enter Max	--	9	0	2
Enter details for P3				
Enter allocation	--	2	1	1
Enter Max	--	2	2	2

Enter details for P4

Enter allocation	--	0	0	2
Enter Max	--	4	3	3

Enter Available Resources	--	3	3	2
---------------------------	----	---	---	---

Enter New Request Details --

Enter pid	--	1		
Enter Request for Resources	--	1	0	2

OUTPUT

P1 is visited(	5	3	2)
P3 is visited(	7	4	3)
P4 is visited(	7	4	5)
P0 is visited(	7	5	5)
P2 is visited(	10	5	7)

SYSTEM IS IN SAFE STATE

The Safe Sequence is -- (P1 P3 P4 P0 P2)

Contd..

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Process	Allocation		Max	Need		
P0	0	1	0	7	5	3
	7	4	3			
P1	3	0	2	3	2	2
	0	2	0			
P2	3	0	2	9	0	2
	6	0	0			
P3	2	1	1	2	2	2
	0	1	1			
P4	0	0	2	4	3	3
	4	3	1			

Programming Code

• Part-1

```
#include<stdio.h>
struct file
{
int all[10];
int max[10];
int need[10];
int flag;
};

int main()
{
struct file f[10];
int fl;
int i, j, k, p, b, n, r, g, cnt=0, id, newr;
int avail[10],seq[10];
printf("Enter number of processes -- ");
scanf("%d",&n);
printf("Enter number of resources -- ");
scanf("%d",&r);
```

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```
for(i=0;i<n;i++)
{
printf("Enter details for P%d",i);
printf("\nEnter allocation\t -- \t");
for(j=0;j<r;j++)
scanf("%d",&f[i].all[j]);
printf("Enter Max\t\t -- \t");
for(j=0;j<r;j++)

scanf("%d",&f[i].max[j]);
f[i].flag=0;
}
printf("\nEnter Available Resources\t -- \t");
for(i=0;i<r;i++)
scanf("%d",&avail[i]);
printf("\nEnter New Request Details -- ");
printf("\nEnter pid \t -- \t");
scanf("%d",&id);
printf("Enter Request for Resources \t -- \t");
```

Part-2

```
for(i=0;i<r;i++)
{
scanf("%d",&newr);
f[id].all[i] += newr;
avail[i]=avail[i] - newr;
}
for(i=0;i<n;i++)
{
for(j=0;j<r;j++)
{
f[i].need[j]=f[i].max[j]-f[i].all[j];
if(f[i].need[j]<0)
f[i].need[j]=0;
}
}
```

```
cnt=0;
fl=0;
while(cnt!=n)
{
g=0;
for(j=0;j<n;j++)
{
if(f[j].flag==0)
{
b=0;
for(p=0;p<r;p++)
{
if(avail[p]>=f[j].need[p])
b=b+1;
else
b=b-1; 2
}
```

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```
}
if(b==r)
{
printf("\nP%d is visited",j);
seq[fl++]=j;
f[j].flag=1;
for(k=0;k<r;k++)
avail[k]=avail[k]+f[j].all[k];
cnt=cnt+1;
printf("(");
for(k=0;k<r;k++)
printf("%3d",avail[k]);
printf(")");
g=1;
}
}
}
if(g==0)
{
printf("\n REQUEST NOT GRANTED -- DEADLOCK
OCCURRED"); printf("\n SYSTEM IS IN UNSAFE STATE");
goto y;
}
}
```

Part-3

```
printf("\nSYSTEM IS IN SAFE STATE");
printf("\nThe Safe Sequence is -- ( );");
for(i=0;i<fl;i++)
printf("P%d ",seq[i]);
printf(")");
y: printf("\nProcess\t\tAllocation\t\tMax\t\t\tNeed\n");

for(i=0;i<n;i++)
{
    printf("P%d\t",i);
    for(j=0;j<r;j++)
        printf("%6d",f[i].all[j]);
    for(j=0;j<r;j++)
        printf("%6d",f[i].max[j]);
    for(j=0;j<r;j++)
        printf("%6d",f[i].need[j]);
    printf("\n");
}
getch();
return 0;
}
```

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